

Ammini College of Engineering

DEPARTMENT OF COMPUTER SCIENCE & ENGINEERING



CS09 407(P)

DATA STRUCTURE LAB

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INTRODUCTION

In computer science, a data structure is a particular way of storing and organizing data in a computer so that it can be used efficiently. Different kinds of data structures are suited to different kinds of applications, and some are highly specialized to certain tasks. For example, B-trees are particularly well-suited for implementation of databases, while compiler implementations usually use hash tables to look up identifiers.

Data structures are used in almost every program or software system. Specific data structures are essential ingredients of many efficient algorithms, and make possible the management of huge amounts of data, such as large databases and internet indexing services. Some formal design methods and programming languages emphasize data structures, rather than algorithms, as the key organizing factor in software design.

Data structures are generally based on the ability of a computer to fetch and store data at any place in its memory, specified by an address — a bit string that can be itself stored in memory and manipulated by the program. Thus the record and array data structures are based on computing the addresses of data items with arithmetic operations; while the linked data structures are based on storing addresses of data items within the structure itself. Many data structures use both principles

STACK

The stack is a very common data structure used in programs. By data structure, we mean something that is meant to hold data and provides certain operations on that data. Stacks hold objects, usually all of the same type. Most stacks support just the simple set of operations we introduced above; and thus, the main property of a stack is that objects go on and come off of the top of the stack.

Here are the minimal operations we'd need for an abstract stack

Push: Places an object on the *top* of the stack.

Pop: Removes an object from the *top* of the stack and produces that object.

IsEmpty: Reports whether the stack is empty or not.

QUEUES

A queue is a first-in first-out data structure (FIFO). Conceptually, the queue data structure behaves like a line. New data is placed at the rear of the queue and when data is removed it is taken from the front. A printer maintains a list of jobs in a queue; the oldest job, the one that has been in the queue the longest, is serviced first.

The operations supported by a queue are as follows:

- void enqueue(Object o) - place object o on the rear of the queue
- Object dequeue() - remove and return the object at the front of the queue
- Object front() - return the object at the front of the queue without removing it
- int size() - return the size of the queue
- boolean isEmpty() - return true if the queue contains no elements, false otherwise

SINGLY-LINKED LISTS

The singly-linked list is the most basic of all the linked data structures. A singly-linked list is simply a sequence of dynamically allocated objects, each of which refers to its successor in the list. Despite this obvious simplicity, there are myriad implementation variations.

DOUBLY-LINKED LIST

A more sophisticated kind of linked list is a doubly-linked list or two-way linked list. Each node has two links: one points to the previous node, or points to a null value or empty list if it is the first node; and one points to the next, or points to a null value or empty list if it is the final node. A doubly-linked list containing three integer values: the value, the link forward to the next node, and the link backward to the previous node.

1. IMPLEMENTATION OF STACK USING ARRAY

Aim

To write a C program to implement stack using array implementation

Algorithm

Step1 : Define the maximum size of stack as 5

Structure

Define a structure stac as follows

Struct stack

```
{  
    int s[size];  
    int top;  
}st;
```

main() function

step 1: start

step 2: Initialize the top of st to -1

step 3: Display the list of choice available for the user to choose
from :

1. Push 2. Pop 3. Display 4.exit

Step 4 : Read the choice from the user as choice

Step 5 : If choice ==1 then read an item to be pushed from user as
item.

If stack is full then display “stack is full” else call the
function push() with item as argument .

step 6 : If choice==2 then check whether “stack is empty”.
If stack is empty then display “stack is empty” else call the Function pop().

Step 7 : If choice ==3 then call the display() function

Step 8 : If choice==4 then go to step9

Step 9 : Stop.

int stfull() – Function to check whether stack is full.

Step 1 : If top of stack st is greater than or equal to size-1 then return 1 else return 0.

void push(int item) – Function to push an item into stack.

Step 1: Increment the value of top of st by 1.

Step 2 : Store the item in the top position of the s as follows:

St.s[st.top]=item

int stempty() – Function to check whether stack is empty.

Step 1 : if top of stack st is -1 then return 1 else return 0.

Int pop() – Function to pop an element from stack.

Step 1: Store the element at top of the st to item

Step 2 : Decrement the value of top by 1

Step 3: Return the value of item.

void display() – Function to display the elements present in the stack

Step 1 : Check whether the stack is empty by calling the function stempty()

If it is empty then display “stack is empty” else go to step2

Step 2 : Display the elements in the stack using for loop shown

below:

```
for(i=st.top;i>=0;i--)  
printf("\n\t%d",st.s[i]);
```

2. IMPLEMENTATION OF QUEUE USING ARRAY

Aim

To write a C program to implement queue using array

Algorithm

Step 1 : Define the maximum size of the queue as 3.

```
#define size 3
```

Main() function

Step 1 : Start

Step 2 : Display the list of choices available to the user to choose

From:

1. Enqueue
2. Dequeue
3. Display
4. Exit

Step 3 : If ch==1, then call the function enqueue() and display()

Step 4 : If ch==2 then call the function dequeue() and display()

Step 5 : If ch==3 then call the function display()

Step 6 : If ch==4 then go to step 5

Step 7 : Stop.

Void enqueue() – function to insert an element in to Queue

Step 1: If rear=size-1 then display “Queue is Full”

Step 2: Increment rear by 1.

Step 3: Read the element from the user and store it in arr[rear]

Step 4: If front==-1 then display “Queue is empty”

Void deueue() – Function to remove an element from queue

- Step 1: If front==-1 then display “Queue is empty”
- Step 2: Dequeue the element at arr[front]
- Step 3: If front==rear then assign -1 to front and rear else increment Front by 1.

Void display()- Function to display the element in queue

- Step 1: If front==-1 then display “Empty Queue”
- Step 2: Display the element in the queue using for loop as follows.
For(i=front;i<=rear;i++)
Printf(“-%d”, arr[i]);

3. SINGLY LINKED LIST

Aim

To write a program to create and perform requested operations on singly linked list.

Algorithm

Self referential structure

Step 1: Define a structure which contains a pointer to a structure of same type as follows:

Structure list :

Integer data – Variable that holds data in node.

Structure list * next – pointer of type struct list next holds the address of the next node in the list.

End Structure

Step 2: Define a variable node which is structure with two members

Struct list

{

Int data;

Struct list *next;

};

typedef struct list node;

Step 3: Make head as the pointer of type struct set the list initially as NULL as given below:

node * head=NULL;

create (int n) – Function for list creation

- Step 1: Declare a pointer variable new of the type node.
- Step 2: Dynamically allocate the memory to the new using malloc function.

```
New=(node *)malloc(size of node);
```

- Step 3: The n is assigned to the data of new
- ```
new → data=n
```

- Step 4: Assign NULL to the next of new
- ```
new → next=NULL
```

- Step 5: new is assigned as the head node of the list.
- ```
Head=new
```

Insbeg(int n) – Adds the node at the beginning of the list.

- Step 1: Declare a pointer variable new of the node.
- Step 2: Dynamically allocate the memory to the new using malloc Function
- ```
New=(node *)malloc(sizeof(node))
```
- Step 3: The n is assigned to the data of new
- ```
New → data=n
```
- Step 4: if head==NULL (i.e., if the list is empty 0 then perform step 5 else goto step 6.

- Step 5:
- 5.1: The new is assigned to the head.
- ```
Head=new
```

- 5.2: Assign NULL to the next of head
- ```
Head → next=NULL
```

Step 6:

6.1: Assign head to the next of new

$\text{New} \rightarrow \text{next} = \text{head}$

6.2: Assign new to head

$\text{Head} = \text{new}$

Insend(int n) – Adds the node at the end.

Step 1: Declare the pointer variables new and temp of type node.

Step 2: Dynamically allocate the memory to new with malloc function

$\text{new} = (\text{node} *) \text{malloc}(\text{size of}(\text{node}));$

Step 3: Assign NULL to the next of new

$\text{New} \rightarrow \text{next} = \text{NULL}$

Step 4: n is assigned to the data of new

$\text{New} \rightarrow \text{data} = \text{n}$

Step 5: Assign the head to temp

$\text{Temp} = \text{head}$

Step 6: Perform step 7 until next of temp is not equal to NULL

$(\text{temp} \rightarrow \text{next} \neq \text{NULL})$  else go to step 8.

Step 7: Assign next of temp to temp

$\text{Temp} = \text{temp} \rightarrow \text{next}$

Step 8: Assign new of the next of temp

$\text{Temp} \rightarrow \text{next} = \text{new}$

Inspos(int n) – Adding node at specified position

Step 1: Declare the pointer variables new and temp of type node.

Step 2: Declare two variables pos, i

Step 3: Dynamically allocate the memory to the new using malloc  
Function

New=(node \*)malloc(sizeof (node));

Step 4: n is assigned to the data of new

New=data=n

Step 5: Read the position from the user (pos)

Step 6: If pos==1 then call the function insbeg() with n as parameter  
Else perform step7

Step 7:

7.1: head is assigned to the temp      temp=head

7.2: for(i=2;i<pos;i++) Perform step 7.3

7.3:

7.3.1: If next of temp is not equal to NULL

(i.e., temp→next!=NULL)perform 7.3.2

7.3.2: next of temp is assigned to temp

Temp=temp→next.

7.4: Assign next of new

New→next=temp→next

7.5: Assign new to the next of temp

Temp→next=new

delbeg()- Deletion of node at beginning.

Step 1: Declare a pointer variable temp of type node.

Step 2: If head==NULL then display “List is empty” else go to step3

Step 3:

3.1: head is assigned to temp.          Temp=head.

3.2: Assign the next of temp to head

Head=temp→next.

Step 4: Deallocate the memory of temp using “true” function

free(temp);

Delend()- Deletion of node at end

Step 1:      Declare two pointer variable s temp and prev of type node

Step 2:      Assign the head to temp.          (temp=head)

Step 3:      Perform step-4 until next of temp is not equal to NULL.  
(i.e., temp→next!=NULL)

Step 4:

4.1: Assign temp to prev          (prev=temp0

4.2: next of temp is assigned to temp.

Temp=temp→next

Step 5:      Assign NULL to the next of prev

Prev→next=NULL

Step 6:      Deallocate the memory of temp using free function

free(temp)

delpos()- Deletion of node at specified position.

Step 1:      Declare two pointer variables temp and prev of type node

Step 2:      Declare two variables pos and i

Step 3:      If head==NULL hen display “List is empty” else go to step 4

Step 4:      Assign head to temp. => temp=head

- Step 5: Read the position for deletion from the user (pos)
- Step 6: If pos==1 then call the function delbeg() else go to step 7
- Step 7:
- 7.1: for(i=2;i<=pos;i++)
- 7.2:
- 7.2.1: if next of temp is not equal to NULL  
(i.e., temp→next!=NULL) perform step 7.2.2
- 7.2.2: Assign the temp to prev (prev=temp)
- 7.3: Assign the next of temp to next of prev  
Prev→next=temp→next
- 7.4: Deallocate the memory of temp using free function  
free(temp);

#### search() – Search for an element

- Step 1: Declare the pointer variable temp of type node
- Step 2: Declare a variable item and initialise i as 1
- Step 3: Read the data to be searched from the user (item)
- Step 4: Assign the head to temp (temp=head)
- Step 5: Perform Step 6 until next of temp is not equal to NULL  
(i.e., temp→next!=NULL) and data of temp is not equal to  
item. (temp→data!=item)
- Step 6:
- 6.1: Assign of temp to temp  
Temp=temp→next
- 6.2: i is incremented by 1

Step 7: If data of temp is equal to item (temp→data==item)  
 Display “Item is found at i” else display “Item is not found”

Display() – Display the list contents.

Step 1: Display a pointer variable temp of type node.

Step 2: Assign head to temp. (temp=head0

Step 3: Display “head→”

Step 4: Perform step -5 until temp is not equal to NULL.

Step 5:

5.1: Display the data of temp

5.2: Assign the next of temp to temp

Temp=temp→next

Step 6: Display “NULL”

Main() function

Step 1: Start

Step 2: Display a list of operation for the user to choose from

1. Single list creation
2. Insert at beginning
3. Insert at end
4. Insert at position
5. Delete at beginning
6. Delete at end
7. Delete at position
8. Search for an element
9. Exit

Step 3: Read the choice entered by the user (choice).

Step 4: If choice==1 then call the function create(no) and display()

Function

- Step 5: If choice==2 then call the function insbeg(no) and display()  
Function
- Step 6: If choice==3 then call the function insend(no) and display()  
Function
- Step 7: If choice==4 then call the function inspos(no) and display()  
Function
- Step 8: If choice==5 then call the function delbeg(no) and display()  
Function
- Step 9: If choice==6 then call the function delbeg(no) and display()  
Function
- Step 10: If choice==7 then call the function delpos(no) and display()  
Function
- Step 11: If choice==8 then call the function search()
- Step 12: If choice==8 then go to step13.
- Step 13: Stop.

## 4. STACK USING LINKED LIST

### Aim

To write a program to perform operation on stack using linked list.

### Algorithm

For linked list operation there is a node having 2 parts. One part is the data part and other pointer part.

Top and t are node type.

### Main() function

- Step 1: Start
- Step 2: Display the menu
- Step 3: If a=1 call the function push then go to step 8  
else go to Step 4
- Step 4: if a=2 call the function pop then go to Step 8 else go to Step 5
- Step 5: if a=3 call the function display then go to Step 8  
else go to Step 6.
- Step 6: If a=4 go to Step 9  
Else go to Step 7
- Step 7: print wrong choice.
- Step 8: if a!=4 go to Step 2.
- Step 9: Stop.

### Function push()

- Step 1: Start
- Step 2:  $t \leftarrow$  new node
- Step 3: Read element in x

- Step 4:  $t \rightarrow \text{data} \leftarrow x$
- Step 5:  $t \rightarrow \text{next} \leftarrow \text{top}$
- Step 6:  $\text{top} \leftarrow t$
- Step 7: print one element is inserted
- Step 8: Return

#### Function Display

- Step 1: Start
- Step 2:  $t \leftarrow \text{top}$
- Step 3: if  $t \neq \text{NULL}$  go to Step 4  
Else go to Step 6
- Step 4: print  $t \leftarrow \text{data}$
- Step 5:  $t \leftarrow t \rightarrow \text{next}$  and go to Step 3
- Step 6: Return

#### Function pop

- Step 1: start
- Step 2:  $t \leftarrow \text{top}$
- Step 3:  $x \leftarrow t \rightarrow \text{data}$
- Step 4:  $\text{top} \leftarrow \text{top} \rightarrow \text{next}$
- Step 5:  $\text{free}(t)$
- Step 6: print the element that is popped.
- Step 7: Return

## 5. QUEUE USING LINKED LIST

### Aim

To write a program to perform operation on queue using linked list.

### Algorithm

For linked list operation there is a node having 2 parts. One is the data part and other is pointer part, pointing to the next node.

Front, rear, t are the node type

#### Main() function

- Step 1: Start
- Step 2: Display the menu
- Step 3: if a=1 call function enqueue then go to Step 8 else gotoStep4
- Step 4: if a=2 call function dequeue then go to Step 8 else go to Step 5
- Step 5: if a=3 call function display then go to Step 8 else goto Step6
- Step 6: if a=4 call function exit then go to Step 9 else go to Step 7
- Step 7: Print wrong choice
- Step 8: if a!=4 go to Step 2
- Step 9: Stop

#### Function enqueue()

- Step 1: Start
- Step 2: Create a new node t
- Step 3: Read the element in x
- Step 4:  $t \rightarrow \text{data} \leftarrow x$
- Step 5: if rear=NULL front $\leftarrow$ t, rear $\leftarrow$ t

Else rear  $\rightarrow$  next  $\leftarrow$  t

Step 6: rear  $\leftarrow$  t

Step 7: rear  $\rightarrow$  next = NULL

Step 8: print one element is inserted

Step 9: Return

#### Function Dequeue()

Step 1: Strat

Step 2: if front = NULL Q is empty return else step 3

Step 3: t  $\rightarrow$  front

Step 4: x  $\leftarrow$  t  $\rightarrow$  data

Step 5: front  $\leftarrow$  front  $\rightarrow$  next

Step 6: free(t)

Step 7: print x

Step 8: Return

## 6. SINGLY LINKED LIST USING POLINOMIAL ADDITION

### Aim

To write a program to perform polynomial addition operation using linked list.

### Algorithm

#### SELF REFRENCIAL DSTRUCTURE

Step 1: Define a structure which contains a pointer to a structure of same type as follows.

#### Structure list

Integer coeff, pow - variables that datas in the node.

Structure list \*next - pointer of type struct list next holds the address of the next node in the list.

End structure

Step 2: Define a new type node which is structure with three members

```
Typedef struct list node;
```

Step 3: Make head as the pointer of type struct.

Set the list initially as NULL

```
Node *head=NULL;
```

Step 4: Repeat the step 3 for poly1,poly2 and poly as given below.

```
Node *poly1=NULL,*poly2=NULL,*poly=NULL;
```

Step 5: Declare the function create(),display() and polyadd() with void as

Return type as return type as given below.

```
Void create(node *);
```

```
Void display(node *);
```

```
Void polyadd(node *,node *,node*);
```

### Main() function

- Step 1: Start
- Step 2: Dynamically allocate the memory to poly1 using malloc  
Function
- Step 3: call the function create() with poly1 as argument
- Step 4: Repeat the steps 2 and 3 for poly2
- Step 5: Display the polynomials poly1 and poly2 by calling the  
Function display() with poly1 and poly2 as arguments  
Respectively.
- Step 6: Dynamically allocate the memory to resultant polynomial  
poly using malloc function
- Step 7: Call the function polyadd() with poly1, poly2 and poly as  
Arguments.

```
Poly(poly1,poly2,poly);
```

- Step 8: Display the resultant polynomial poly by calling the  
display() function
- Step 9: Stop.

### Create(node \*p)

- Step 1: Declare a character variable ch
- Step 2: Read the value for coefficient from the user and store it in  
Coeff of p as follows

```
Scanf(“%d”,&p→coeff);
```

- Step 3: Read the value for power from the user and store it in pow of

p as follows.

```
scanf("%d",&p->pow);
```

Step 4: If the user wants to continue then perform Step 5  
else go to Step 6

Step 5: 5.1: Dynamically allocate the memory to next of p using  
malloc() function

```
p->next=(node *)malloc(size of (node));
```

5.2: next of p is assigned to p

```
p=p->next;
```

Step 6: Assign NULL to the next of p

```
p->next=NULL;
```

display(node \*p)

Step 1: Perform the Step 2 until p is not equal to NULL

```
p!=NULL
```

Step 2: 2.1: Display the coefficient and power of the each term in  
The polynomial with the help of following statement.

```
printf("%d%d",p->coeff,p->pow);
```

2.2: next of p is assigned to p

```
p=p->next
```

2.3: If p is not equal to NULL (i.e., p!=NULL) then

```
display '+'
```

void polyadd(node \*p1,node \*p2,node \*p)

Step 1: Perform Step 2 to Step 9 until p1 not equal to NULL and p2  
not equal to NULL

(i.e.,  $p1 \neq \text{NULL} \&\& p2 \neq \text{NULL}$ )

Step 2: If pow of p2 equal to pow of p1 (i.e.,  $p1 \rightarrow \text{pow} == p2 \rightarrow \text{pow}$ ) then perform Step 3 else go to Step 4.

Step 3: 3.1: Assign pow of p1 to pow of p

$p \rightarrow \text{pow} = p1 \rightarrow \text{pow}$

3.2: Assign the sum of the coeff of p1 and p2 to coeff of p

$p \rightarrow \text{coeff} = p1 \rightarrow \text{coeff} + p2 \rightarrow \text{coeff};$

3.3: Assign next of p1 to p1. ( $p1 = p1 \rightarrow \text{next}$ )

3.4: Assign next of p2 to p2. ( $p2 = p2 \rightarrow \text{next}$ )

Step 4: If pow of p1 is greater than the pow of p2 (i.e.,  $p1 \rightarrow \text{pow} > p2 \rightarrow \text{pow}$ ) then perform Step 5 else go to Step 6

Step 5: 5.1: pow of p1 is assigned to the pow of p

$p \rightarrow \text{pow} = p1 \rightarrow \text{pow}$

5.2: coeff of p1 is assigned to the coeff of p

$p \rightarrow \text{coeff} = p1 \rightarrow \text{coeff}$

5.3: Assign next of p1 to p1

$P1 = p1 \rightarrow \text{next}$

Step 6: 6.1: Assign pow of p2 to pow of p

$p \rightarrow \text{pow} = p2 \rightarrow \text{pow};$

6.2: Assign coeff of p2 to coeff of p

$p \rightarrow \text{coeff} = p2 \rightarrow \text{coeff};$

6.3: Assign next of p2 to p2

P2=p2→next;

Step 7: If p1 is not equal to NULL or p2 is not equal to NULL  
(i.e., p1!=NULL||p2!=NULL) then perform Step 8  
else go to Step 9

Step 8: 8.1: Dynamically allocate the memory to next of p using  
malloc() function.

p→next=(node \*)malloc(sizeof(node));

8.2: Assign next of p to p

P=p→next;

Step 9: Assign NULL to next of p  
p→next=NULL;

Step 10: Perform Steps 11 to Step 17 until p1 is not equal to NULL  
or p2 is not equal to NULL.  
(i.e., p1!=NULL||p2!=NULL)

Step 11: If p1 not equal to NULL then perform Step 12.  
Else go to Step 13

Step 12: 12.1: Assign pow of p1 to pow of p

p→pow=p1→pow

12.2: Assign coeff of p1 to coeff of p

p→coeff=p1→coeff

12.3: Assign next of p1 to p1 (p1=p1→next)

Step 13: If p2 is not equal to NULL then perform Step 14  
else go to Step 15

Step 14:      14.1: Assign pow of p2 to pow of p  
                   $p \rightarrow \text{pow} = p2 \rightarrow \text{pow}$

                  14.2: Assign coeff of p2 to coeff of p  
                   $p \rightarrow \text{coeff} = p2 \rightarrow \text{coeff}$

                  14.3: Assign next of p2 to p2  
                   $P2 = p2 \rightarrow \text{next}$

Step 15:      If p1 not equal to NULL or p2 not equal to NULL  
                  (i.e.,  $p1 \neq \text{NULL} \parallel p2 \neq \text{NULL}$ ) then perform step 16  
                  Else go to Step 17.

Step 16:      16.1: Dynamically allocate the memory to next of p using  
                  Malloc() function  
                   $p \rightarrow \text{next} = (\text{node} *) \text{malloc}(\text{sizeof}(\text{node}));$

                  16.2: Assign next of p to p  
                   $P = p \rightarrow \text{next};$

Step 17:      Assign NULL to next of p  
                   $p \rightarrow \text{next} = \text{NULL}$

## 7. DOUBLY LINKED LIST

### Aim

To write a program to create and perform various operations, such as insertion, deletion and searching for an element on doubly linked list.

### Algorithm

#### SELF- REFERENTIAL STRUCTURE

Step 1: Define a structure which contains a pointer to a structure of same as follows:

Structure list

Integer data – Variable that holds data in node.

Structure list \* next, \*prev- pointer of type struct list next holds the address of the next node in the list.

End Structure

Step 2: Define a variable node which is structure with two members

Struct list

{

Int data;

Struct list \*prev, \*next;

};

typedef struct list node;

Step 3: Make head as the pointer of type struct. Set the list initially as NULL as given below.

Node \*head=NULL

Create(int n) - Function for list creation.

- Step 1: Create a pointer variable new of type node
- Step 2: Dynamically allocate the memory to new using malloc()  
New=(node \*)malloc(sizeof(node));
- Step 3: Assign n to the data of new  
New→data=n
- Step 4: Assign NULL to the next of new  
New→next=NULL
- Step 5: Assign NULL to the prev of new  
New→prev=NULL
- Step 6: Assign new to head  
Head=new

Void insbeg(int n) - Function for insertion at beginning

- Step 1: Declare a pointer variable new of type node.
- Step 2: Dynamically allocate the memory to new using malloc()  
New= (node \*)malloc(sizeof node));
- Step 3: Assign n to the data of new  
New→data=n
- Step 4: If head==NULL then perform Step 5 else go to Step 6
- Step 5:
- 5.1: Assign new to head head=new
- 5.2: Assign NULL to the prev of new  
New→prev=NULL
- 5.3: Assign NULL to the nextof new  
New→next=NULL

Step 6:

6.1: Assign head to the next of new

New→next=head

6.2: Assign new to the prev of head

Head→prev=new

6.3: Assign new to head      head=new

6.4: Assign NULL to prev of head

Head→prev=NULL

Void insend(int n) - Function to perform insertion at end

Step 1: Declare two pointer variables new and temp of type node

Step 2: Dynamically allocate the memory to new using malloc()

New= (node \*)malloc(sizeof (node));

Step 3: Assign n to the data of new

New→data=n

Step 4: Assign NULL to the next of new

New→next=NULL

Step 5: Assign head to temp

Temp=head

Step 6: Perform Step 7 until next of temp is not equal to NULL

(i.e., temp→next!=NULL)

Step 7: next of temp is assigned to temp

Temp=temp→next

Step 8: Assign new to the next of temp

Temp→next=new

Step 9: Assign temp to the prev of new

New → prev = temp

Void inspos(int n)- Function to insert at given position

Step 1: Declare a variable x of integer type.

Step 2: Declare a pointer variable new, temp and temp1 of type node.

Step 3: Dynamically allocate the memory to new using malloc()

Step 4: Assign n to the data of new

New → data = n;

Step 5: Read the position from the user (x)

Step 6: If x == 1 then call the function insbeg(n) else go to Step 7

Step 7: 7.1: Assign head to temp      temp = head

7.2: for(i=2; i < x; i++)      perform Step 7.3

7.3:

7.3.1 : If next of temp is not equal to NULL

(i.e., temp → next != NULL) perform Step 7.3.2

7.3.2 : next of temp is assigned to temp1

Step 8: next of temp is to temp1

Temp1 = temp → next

Step 9: temp is assigned to prev of new

New → prev = temp

Step 10: temp1 is Assigned to next of new

New → next = temp1

Step 11: new is assigned to next of temp

Temp → next = new

Step 12: new is assigned to prev of temp1

Temp1 → prev=new

Void disbeg() – function to delete at beginning

Step 1: Declare a pointer variable temp of type node

Step 2: If head==NULL then display “ List is empty”  
else go to Step 3

Step 3:

3.1: head is assigned to temp =>temp=head

3.2: next of head is assigned to head

Head=head→next

3.3: Assign NULL to the prev of head

Head→prev=NULL

Step 4:

4.1: Perform Step 4.2 until temp→next!=NULL

4.2: next of temp is assigned to temp

Temp=temp→next

4.3: prev of temp is assigned to temp1

Temp1=temp→prev

4.4: Assign NULL to prev of temp

Temp→prev=NULL

Step 5: Assign NULL to the next of temp1

Temp1→next=NULL

Step 6: Deallocate the memory of temp using free() function

Free(temp);

Void delpos() – Function to delete at a given position

Step 1: Declare pointer variables temp, temp1, temp2 of type node.

Step 2: Declare two variables i and x of integer type

Step 3: if head==NULL then display “List is empty”

Else go to Step 4

Step 4:

4.1: Assign head to temp                   =>temp=head

4.2: Read the position from the user (x)

4.3: If x==1 then call the function delbeg()

Else perform Step 4.4

4.4:

4.4.1: for(i=2;i<=x;i++) perform steps 4.4.2 to 4.4.5

4.4.2: If temp→next!=NULL perform Step 4.4.3

Else go to Step 4.4.4

4.4.3: Assign next of temp to temp

Temp=temp→next

4.4.4: Assign prev of temp to temp1

Temp1=temp→prev

4.4.5: Assign next of temp to temp2

Temp2=temp→next

4.4.6 :Assign temp1 to prev of temp2

Temp2→prev=temp1

4.4.7: Assign temp2 to next of temp1

Temp→next=temp2

4.4.8: Assign NULL to the next of temp

Temp→next=NULL

4.4.9: Assign NULL to the prev of temp

Temp→prev=NULL

Step 5: Deallocate the memory of temp using free()

Free(temp0;

Void search()- Function to search for an element

Step 1: Declare a pointer variable temp of the type node.

Step 2: Declare two variables item of i of integer type.

Step 3: Read the data to be searched from the user

Step 4: Assign head to temp      temp=head

Step 5: Perform Step 6 until  
(temp→next!=NULL &&temp->data!=item)

Else go to Step 7

Step 6:

6.1: Assign next of temp to temp

Temp=temp→next

6.2: Increment the variable i by 1

Step 7: If (temp→data==item) display “Item is found at i”

Else go to Step 8

Step 8: Display “Item not found”

Void display()- Function to display the list

Step 1: declare a pointer variable temp of type node

Step 2: Assign head to temp      temp=head

- Step 3: Display head
- Step 4: Perform step 5 until temp!=NULL  
Else go to Step 6
- Step 5:
- 5.1: Display the data
  - 5.2: Assign next temp to temp  
Temp=temp→next
- Step 6: Display “NULL”

### Main() Function

- Step 1: Start
- Step 2: Display a list of operations for the user to choose from
- 1. Single list creation
  - 2. Insert at beginning
  - 3. Insert at end
  - 4. Insert at position
  - 5. Delete at beginning
  - 6. Delete at end
  - 7. Delete at position
  - 8. Search for an element
  - 9. Exit
- Step 3: Read the choice from the user
- Step 4: If choice=1 call the function create(no) and display()
- Step 5: If choice=2 call the function insbeg(no) and display()
- Step 6: If choice=3 call the function insend(no) and display()
- Step 7: If choice=4 call the function inspos(no) and display()
- Step 8: If choice=5 call the function delbeg() and display()
- Step 9: If choice=6 call the function delend() and display()

Step 10: If choice=7 call the function delpos() and display()

Step 11: If choice=8 call the function search()

Step 12: If choice=9 then go to Step 13

Step 13: Stop.

## 8. QUICK SORT

### Aim

To sort n number in ascending order using quick sort algorithm

### Algorithm

- Step 1: Start
- Step 2: Read n numbers to be sorted in to an array
- Step 3: Call the function quick sort passing variable
- Step 4: print the numbers
- Step 5: stop

### Function Quick sort()

Initial state  $m \leftarrow 0, n \leftarrow n-1$

- Step 1: Start
- Step 2: If  $m < n$  go to Step 3
- Step 3:  $i \leftarrow m, j \leftarrow n, p \leftarrow a[m]$
- Step 4:  $i \leftarrow i+1$
- Step 5: if  $((a[j] \geq p) \& \& (j > m))$  go to Step 7  
Else go to Step 8
- Step 6: if  $((a[i] \leq p) \& \& (j \leq n))$  go to Step 4  
Else go to Step 6
- Step 7:  $j \leftarrow j-1$  go to Step 6
- Step 8: if  $i < j$  go to Step 9  
Else go to Step 10
- Step 9: swap  $a[i] \& a[j]$
- Step 10: if  $i \leq j$  go to Step 4

Else go to Step 11

Step 11: swap  $a[m]$  &  $a[j]$

Step 12: Call function quick sort by passing  $a, m, j-1$

Step 13: Call function quick sort by passing  $a, j+1, n$

Step 14: Return.



- Step 4: call function merge and pass  $x, y, i, i+1-1, i+2l-1$
- Step 5:  $i=l+2*1$  go to Step 3
- Step 6: if  $(i+i-1) < n$  do step 7  
Else go to Step 8
- Step 7: call function & pass  $x, y, i, i+1-1, n-1$  & go to Step 12
- Step 8:  $k=j$
- Step 9: if  $k < n$  do step 10  
Else go to Step 12
- Step 10:  $y[k]=x[k]$
- Step 11:  $i=k+1$  & go to Step 9
- Step 12: Return

### Function Merge

- Step 1: Start
- Step 2:  $i=l, j=m+1, k=l$
- Step 3: if  $i \leq m \& j \leq n$  do Step 4  
Else do Step 8
- Step 4: if  $x[i] < x[j]$  do Step 5  
Else do Step 6
- Step 5:  $y[k]=x[j], i=j+1$  go to Step 7
- Step 6:  $y[k]=x[j], j=j+1$
- Step 7:  $k=k+1$  go to Step 3
- Step 8: if  $i > m$  do Step 9 else go to Step 11
- Step 9: if  $j \leq n$  do Step 11  
Else go to Step 13

Step 10:  $y[k]=x[j], j=j+1, k=k+1$  go to Step 7

Step 11: if  $i \leq m$  do Step 12 else go to Step 13

Step 12:  $y[k]=x[i], i=i+1 k=k+1$  go to Step 11

Step 13: Return

## 10. INFIX TO POSTFIX CONVERSION

### Aim

To write a program to convert a given infix expression to postfix expression

### Algorithm

#### Main() Function

- Step 1: Start
- Step 2: Read the maximum number of character from the user
- Step 3: for(i=0;i<max;i++), read the characters from the user and store in infix[i] and if infix[i]=='#' then go to Step 4
- Step 4: for(i=0;i<max;i++) display the infix expression  
                                           Puchar(infix[i])
- Step 5: Call the function infix-postfix()
- Step 6: for(i=0;i<max;i++) display the postfix expression  
                                           Puchar(postfix[i])
- Step 7: stop

Infix-postfix() - function to convert an infix expression to postfix

#### Expression

- Step 1: Declare a counter variable i, a character variable next and initialise p as 0
- Step 2: stack[top]='#'
- Step 3: len=strlen(infix)
- Step 4: infix[len]='#'
- Step 5: for(i=0;infix[i]!='#';i++) perform 6 to 11

Step 6: if infix[i]='(' then call push() with infix[i] as argument

Step 7: if infix[i]=')' then perform Step 8

Step 8

8.1: next=pop()

8.2: postfix[p++]=next;

Step 9: if infix[i]='+', '-', '\*', '/', '%', or '^' then perform Step 10

Step 10:

10.1: precedence=prec(infix[i])

10.2: Perform Step 10.3 until the operator!='#' and

Precedence of the operator <=store in stak[]

10.3: postfix[p++]=pop()

10.4: call push() with infix[i] as argument

Step 11: if infix[i]!=anyone mentioned above store the character in the postfix[]

Step 12: perform Step 13 until the operator in stack[]!='#'

Step 13:

13.1: postfix[p++]=pop()

13.2: postfix[p]='\0'

Void push(char symbol)- Function to perform push operation in stack

Step 1: If (top>max) then display "Stack overflow"

Else go to Step 2

Step 2:

2.1: top++

2.2: stak[top]=symbol

Int pop()- Function to perform pop operation in stack

Step 1: If top==-1 then display “Stack underflow”

Else go to Step 2

Step 2: Return the operator stored in stack[]

Return (stack[top--])

Int prec(char symbol)- Function to assign precedence to an operator

Step 1: If symbol=='(' then return 0

Step 2: If symbol=='+' or '-' then return 1

Step 3: If symbol=='\*' or '/' or '%' then return 2

Step 4: If symbol=='^' then return 3

Step 5: Return 0

## 11. IMPLEMENTATION OF EXPRESSION TREE CONSTRUCTION AND TREE TRAVERSAL

### Aim

To write a c program to create an expression tree from the given postfix expression and find the prefix and infix expression for that expression

### Algorithm

Self referential structure

Step 1: Define a structure which contains a pointer to a structure

```
Struct tree
{
 Char data;
 Struct tree *lchild;
 Struct tree *rchild;
};
```

Step 2: Define a variable node for that structure using typedef

```
Typedef struct tree node;
```

Step 3: Declare an array pointer of size 50 for the node

```
Node * stack[50];
```

Void push(node \*t)- Function to push an element into stack

Step 1: top++

Step 2: stack[top]=t

Node pop() - function to pop out an element from stack

Step 1: Declare pointer variable t of type node

```
Node *t;
```

Step 2: Assign the element at stack[top--] to t and return it.

T=stack[top--]

Return t;

void inorder(node \*ptr)- Function to display inorder expression

Step 1: If ptr!=NULL then perform Step 2 to 4

Step 2: Call the recursive function inorder() with ptr→lchild as Argument

Step 3: Print the data at ptr

Step 4: Call inorder() with ptr→rchild as argument recursively

void preorder(node \*ptr) – Function to display preorder expression

Step 1: If ptr!=NULL then perform steps 2 to 4

Step 2: Print the data at ptr

Step 3: Call preorder() with ptr→lchild as argument, recursively

Step 4: Call preorder() with ptr→rchild as argument, recursively

void postorder(node \*ptr)- Function to display the postorder expression

Step 1: if ptr!=NULL then Steps 2 to 4

Step 2: Call postorder() with ptr→lchild as argument recursively

Step 3: Call postorder() with ptr→rchild as argument recursively

Step 4: Print the data at ptr

node \*getnode(char s)- Function to store the operators in the tree

Step 1: Declare a pointer variable t of type node

Step 2: Dynamically allocate the memory to t using malloc()

T=(node \*0malloc(sizeof(node));

Step 3: Store the s at the data of t

Step 4: Assign the s at the data of t

Step 5: Return t

void logicalview(node \*t,int col,int row,int wid,int call)

Step 1: call gotoxy() with col and row as arguments

Step 2: if t!=NULL then perform Step 3 to 7

Step 3: if call==0 then display"Root"

Step 4: Call the function gotoxy with col and row+1

Step 5: Display data of t.

Step 6: Call logicalview() recursively with t-→rchild,col+wid,  
Row+2,wid/2 and call+1 as arguments

Main() function

Step 1: Start

Step 2: Declare 2 pointer variables type node

Node \*temp,\*new1

Step 3: Read the positive expression from the user as str[50]

Step 4: for(i=0;str[i]!='\0';i++)

Step 5: if the character at str[i] is alphanumeric then perform Step 6  
Else go to Step 7

Step 6

6.1: call getnode(str[i]) and assign it to new

6.2: call the push() with new1 as argument

Step 7:

7.1: call getnode(str[i]) and assign it to new1

7.2: Assign new1 to temp

7.3: call pop() for left and right childs of temp

```
Temp→rchild=pop();
```

```
Temp→lchild=pop();
```

7.4: call push() with temp as argument

Step 8: Display expression tree by calling logicalview() with temp

Step 9: Display the preorder, inorder and postorder expression by calling preorder(),inorder() and postorder() with temp as argument respectively.

Step 10: Stop

## 12. IMPLEMENTATION OF BINARY SEARCH TREE

### Aim

To write a C program to create a binary search tree and perform various operations like insertion, deletion, find an element, find minimum element, maximum element and display the tree on it.

### Algorithm

Self referential structure

Step 1: Define a structure which contains a pointer to the structure of same type

```
struct tree
{
 Int data;
 Struct tree *left;
 Struct tree *right;
};
```

Step 2: Define a variable node of the tree

```
Typedef struct tree node;
```

Step 3: Declare a pointer variable t for node and assign it to NULL

```
Node *t=NULL;
```

Step 4: Declare a global variable call of integer type and

Assign it to 0

```
Int call=0;
```

Node \* makeempty(node \*)- Function for deleting all the element in the tree

- Step 1: If  $t \neq \text{NULL}$  then perform Step 2 to 4
- Step 2: Call the function `makeempty()` recursively with  $t \rightarrow \text{left}$  as argument to delete left child of a node
- Step 3: Call the function `makeempty()` recursively with  $t \rightarrow \text{right}$  As argument to delete right child of a node.
- Step 4: Deallocate the memory of  $t$  using `free()`function
- Step 5: Return `NULL`

Node \*find (int x,node \*t)- Function for finding the element in tree

- Step 1: If  $t == \text{NULL}$  then return `NULL` else go to Step 2.
- Step 2: If  $x < t \rightarrow \text{data}$  then return `find(x,t → left)`  
Else go to Step 3
- Step 3: If  $x > t \rightarrow \text{data}$  then return `find(x,t → right)`

Node \*findmin(node \*t)- Function for finding minimum element.

- Step 1: If  $t == \text{NULL}$  then return `NULL` else go to Step 2
- Step 2: If there is no left child then return the element at  $t$  else  
Call the function `findmin(t → left)` recursively

Node \*findmax(node \*t)- Function for finding maximum element

- Step 1: If  $t == \text{NULL}$  then return `NULL` else go to Step 2
- Step 2: If there is no right child then return the element at  $t$ .  
Else call recursive function `findmax[t → right]`

Node \*insert(int x,node \*t)- Function to insert element

- Step 1: If  $t == \text{NULL}$  perform Step 2 else go to Step 3
- Step 2:

2.1: Dynamically allocate the memory to t using malloc()

$T=(\text{node } *)\text{malloc}(\text{sizeof}(\text{node}));$

2.2: Store x in data of t

2.3: Make the left and right child of t to point to NULL

Step 3:

3.1: If  $x < t \rightarrow \text{data}$  then insert x as left child of t

3.2: If  $x > t \rightarrow \text{data}$  then insert x as right child of t

Step 4: Return t

Node \* delete (int x, node \*t)- Function to delete element

Step 1: Declare a pointer variable temp of type node.

Step 2: If  $t == \text{NULL}$  then display "Element not found"

Else go to Step 3

Step 3: If  $x < t \rightarrow \text{data}$  then traverse along left child for deleting x by calling the function  $\text{delete}(x, t \rightarrow \text{left})$  and Assign it to  $t \rightarrow \text{left}$  else go to Step 4.

Step 4: If  $x > t \rightarrow \text{data}$  then traverse along right child for deleting x by calling  $\text{delete}(x, t \rightarrow \text{right})$  and Assign it to  $t \rightarrow \text{right}$

else go to Step 5

Step 5: If both left and right child exist for t then perform Step 6

Else go to Step 7

Step 6:

6.1: Call  $\text{findmin}(t \rightarrow \text{right})$  and assign it to temp

$\text{Temp}=\text{findmin}(t \rightarrow \text{right});$

6.2: Assign data of temp to data of t

$t \rightarrow \text{data}=\text{temp} \rightarrow \text{data};$

6.3: call delete( $t \rightarrow \text{data}, t \rightarrow \text{right}$ ) and assign it to  $t \rightarrow \text{right}$   
 $t \rightarrow \text{right} = \text{delete}(t \rightarrow \text{data}, t \rightarrow \text{right});$

Step 7:

7.1: Assign t to temp

7.2: If  $t \rightarrow \text{left} == \text{NULL}$  then Assign the data at right child  
 To t, else go to Step 7.3

7.3: If  $t \rightarrow \text{right} == \text{NULL}$  then assign  
 the data at left child of t to t.

Step 8: Return t

Void display(node \*t, int col, int row, int wid, int call)

Step 1: Call gotoxy() with col and row as argument

Step 2: If  $t != \text{NULL}$  then perform Step 3 to 7

Step 3: If  $\text{call} == 0$  then display ROOT

Step 4: Call gotoxy() with col and row+1 as arguments

Step 5: Display data of t

Step 6: call display( $t \rightarrow \text{left}, \text{col} - \text{wid}, \text{row} + 2, \text{wid} + 1, \text{call} + 1$ );

Step 7: call display( $t \rightarrow \text{right}, \text{col} + \text{wid}, \text{row} + 2, \text{wid} / 2, \text{call} + 1$ );

Main() function

Step 1: Start

Step 2: Display list of operations for the user to choose from

1. Insert
2. Delete
3. Find
4. Findmin
5. Findmax
6. Display
7. Makeempty

## 8. Exit

- Step 3: Read the choice from the user choice
- Step 4: If choice==1 then read the data to be inserted as c and call the functions insert(c,t) and display (t,40,2,16,call)
- Step 5: If choice==2 then read the data to be deleted as c  
If t==NULL, then display “tree is empty”  
Else call delete(c,t) and display (t,40,2,16,call)
- Step 6: If choice==3 then read data to be searched as c  
If t==NULL then display “tree is empty”  
Else call find(c,t)
- Step 7: If choice==4 then check whether tree is empty or not.  
If tree is not empty then call findmin() and Assign it to r
- Step 8: If choice==5 then check whether tree is empty or not.  
If tree is not empty then call findmax(t) and Assign to r
- Step 9: If choice==6 then check whether tree is empty or not.  
If tree is not empty then call display(t,40,2,16,call)
- Step 10: If choice==7 then call makeempty(t) and assign it to t.  
Display “Tree is empty”
- Step 11: If choice==7 then go to Step 12
- Step 12: Stop.

### 13. IMPLEMENTATION OF HASHING TECHNIQUE USING LINEAR PROBING

#### Aim

To create a hash table based on modulo division method using linear probing method, in C program

#### Algorithm

Step 1: Define max as 10 as shown below

```
#define max 10
```

Create(int num) – Function to create and perform hash function

Step 1: Assign the  $\text{num} \% 10$  to  $\text{key} \Rightarrow \text{key} = \text{num} \% 10$

Step 2: Return the value of key

Iprobe(int a[max],int key, int num)- Function to insert an element in to  
hash table

Step 1: Initialise flag to 0

Step 2: If  $\text{a}[\text{key}] == -1$  then store num at  $\text{a}[\text{key}]$  position in the array  
else go to Step 3

Step 3: Perform Step 4 until  $i < \text{max}$

Step 4: Find how many position are filled in the table using the  
Condition

```
If($\text{a}[\text{i}] \neq -1$)
```

```
Count++;
```

```
i++;
```

Step 5: If  $\text{count} == \text{max}$  then display “Hash table is full”

Step 6: for( $\text{i} = \text{key} + 1; \text{i} < \text{max}; \text{i}++$ ) perform Step 7

Step 7: If  $a[i] == -1$  perform Step 8

Step 8:

8.1: Assign num to  $a[i]$

8.2: Assign 1 to Flag and break the loop

Step 9: If  $i == \text{max}$  perform Step 10

Step 10: Initialise  $i$  to 0 and perform Step 8 until  $i < \text{max}$

#### Void display(int a[max])

Step 1: Display the contents in the hash table using for loop as follows:

```
for(i=0;i,max;i++)
printf("\n\t\t%d\t\t%d",i,a[i]);
```

#### main() function

Step 1: Start

Step 2: Intialise all elements of hashtable as -1 by using for loop as

```
For(i=0;i<max;i++)
```

```
a[i]=-1;
```

Step 3: Read the number from the user as num

Step 4: Call the function create() with num as argument and assign it to key

Step 5: Call the function lprobe() with a, key and num as arguments

Step 6: Call the function disply() with a as argument

Step 7: Perform steps 3,4,5 and 6 until user wants to continue

Step 8: Stop

## 14. IMPLEMENTATION OF DIJKSTRA'S ALGORITHM

### Aim

To write a C program to find minimum distance from source vertex to all other vertices using Dijkstra's algorithm

### Algorithm

Structure

Step 1: Define a structure edge as follows

Struct edge

{

    Int vs,vd,wt;

    }a[20];

    Vs- Source vertex

    Vd-Destination vertex

    Wt- Weightage

Void di(int n)- Function to find shortest distance from single node to all nodes.

Step 1: Make all the vertices unvisited using the for loop

For(i=1;i<=n;i++)

Str[j]=0;

Step 2: Create edges for each non-zero element in the given adjacency matrix. Edge should be created with source vertex vs, destination vertex vd, and weightage wt.

Step 3: Read the starting vertex from the user as s and make it visited

St[s]=1

Let the predecessor of s as 1 and the distance be 0

Step 4: Make the other nodes distance be 9999 and their predecessor as 0

```
for(i=0;i<=n;i++)
{
 If(i!=s)
 {
 Pred[i]=0;
 Dist[i]=9999;
 }
}
```

Step 5: Find the distance from the vertex which is visited to its adjacent vertex using the for loop

```
for(i=1;i<=k;i++)
{
 If((st[a[i].vs]==1)&&(st[a[i].vd]==0))
 {
 M[+t]=(a[i].wt+dist[a[i].vs]);
 }
}
```

Step 6: Consider the vertex with smallest distance as visited and its predecessor be vs and its distance be shortest distance between vs and vd

```
for(i=1;i<=k;i++)
```

```
{
 If((st[a[i].vs]==1)&&(st[a[i].vd]==0))
 {
 If((a[i].wt+dist[a[i].vs])==sm)
 {
 St[a[i].vd]=1;
 Pred[a[i].vd]=a[i].vs;
 Dist[a[i].vd]=a[i].wt+dist[a[i].vs];
 }
 }
}
```

Step 7: Repeat steps 4,5 and 6 until all nodes are visited

Step 8: Print the predecessor value of all nodes and their corresponding distance values.

#### Main() function

Step 1: Start

Step 2: Read no.of vertices from the user as n

Step 3: Read the adjacency matrix from the user as g[i][j]

Step 4: Call the function dj() with n as argument

Step 5: Stop

## 15. DFS & BFS

### Aim

Write an algorithm and corresponding program for depth first search and breadth first search

### Algorithm

- Step 1: DES(v) - v is the vertex
- Step 2: mark[v]:visited
- Step 3: for each vertex w on L[u] do
- Step 4: if mark – reverse call to DFS
- Step 5: DFS[w] – reverse call to DFS

BFS search for each vertex v that are visited

- Step 1: BSF(v)
- Step 2: Mark[v]
- Step 3: equal(v,q)
- Step 4: Repeat step while not empty(a) do
- Step 5: x=front(q)
- Step 6: dequeue(q)
- Step 7: for each vertex adjacent to x do
- Step 8: if mark[y]=unvisited then
- Step 9: mark[s]=visited
- Step 10: enqueue(y,q)
- Step 11: insert((x,y),1)